



max planck institut  
informatik

# Hamilton cycles in the random geometric graph

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**Montpellier, June 21, 2012**

# Random graphs and related models

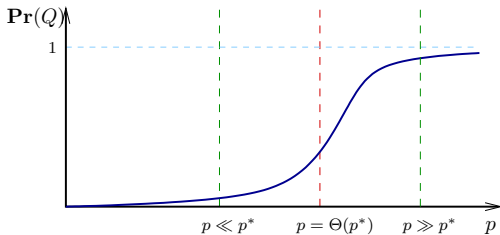
- Erdős-Rényi models:  $\mathcal{G}(n, p)$  and  $\mathcal{G}(n, m)$
- Variations: fixed degree sequence, regular graphs...
- Power-law degree sequence: inhomogeneous random graphs, preferential attachment, internet graph...
- Random boolean formulas
- Statistical mechanics: Ising model, Potts model
- Proximity graphs: **Random geometric graphs**, nearest neighbour graphs, Delaunay graphs, models of wireless networks...



# Threshold functions

## Usual setting:

- $n \rightarrow \infty$  vertices,
- $p = p(n)$  “density” parameter,
- property  $Q$



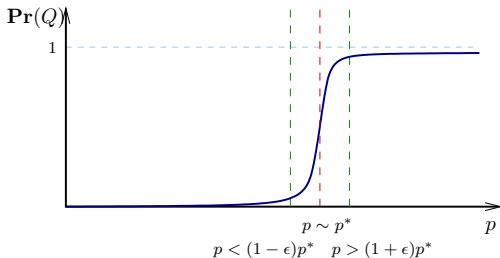
Example: existence of triangles



# Threshold functions (sharp)

## Usual setting:

- $n \rightarrow \infty$  vertices,
- $p = p(n)$  “density” parameter,
- property  $Q$



Example: giant component, connectedness, Hamilton cycles...

# Random graph process

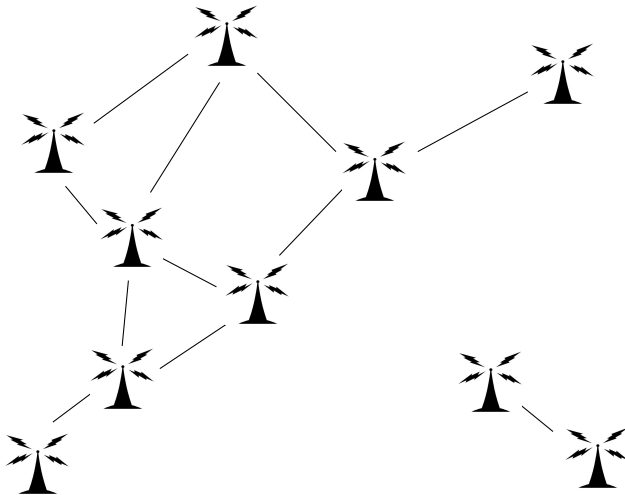
- Start with empty graph on  $n$  vertices; add edges one by one; end with the complete graph

$$\mathcal{G}_0, \mathcal{G}_1, \dots, \mathcal{G}_m, \dots, \mathcal{G}_{\binom{n}{2}}$$

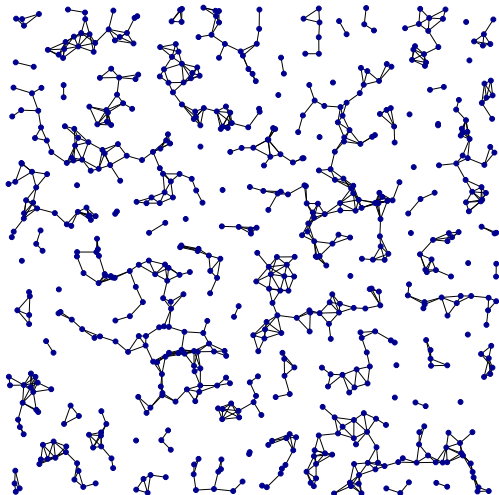
- $\mathcal{G}_m$  is distributed like  $\mathcal{G}(n, m)$
- We look for “hitting time” properties



# Wireless networks



# Random geometric graph



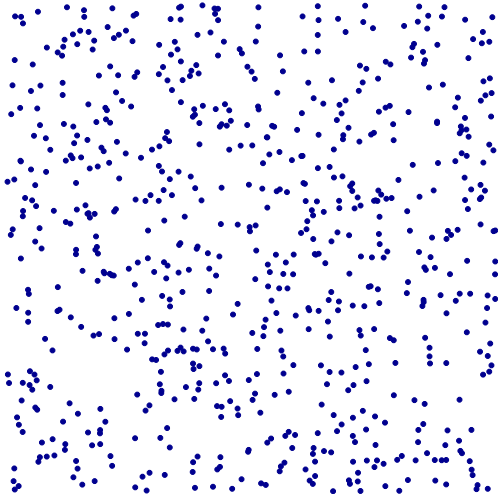
(Gilbert 1961)

$n$  vertices

radius  $r = r(n)$

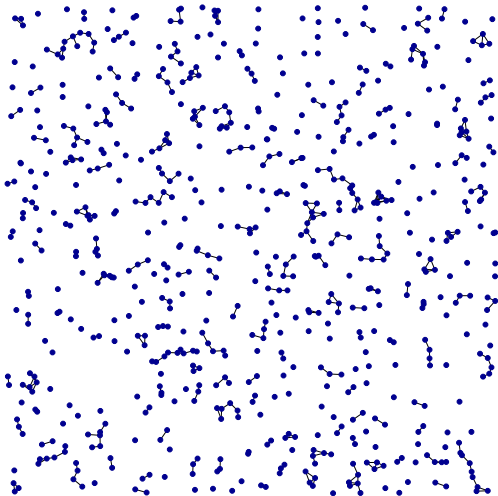
$n \rightarrow \infty$

Random process:  $0 \leq r \leq \sqrt{2}$

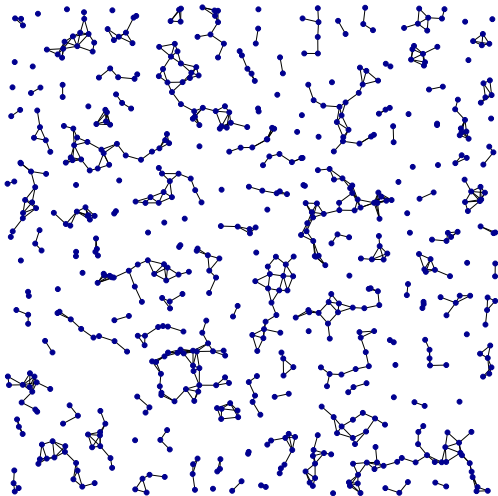




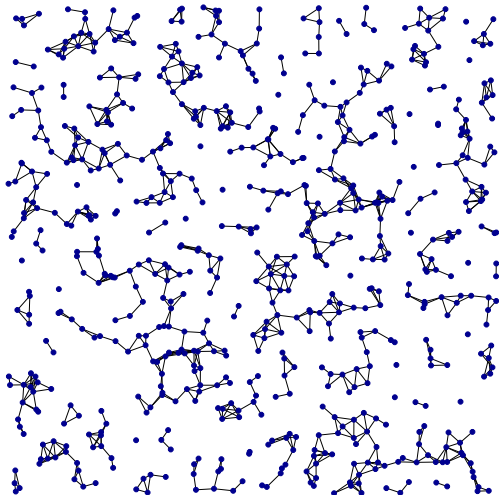
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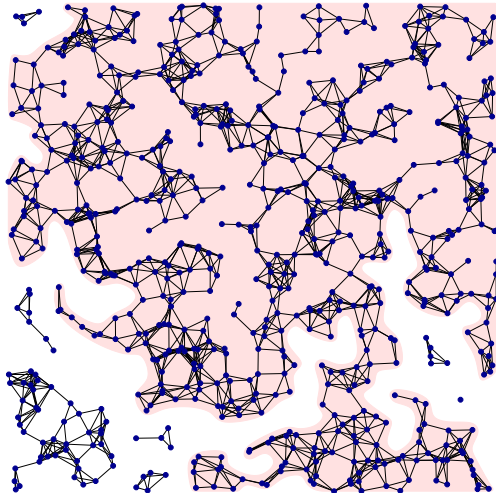


Random process:  $0 \leq r \leq \sqrt{2}$



no giant component yet

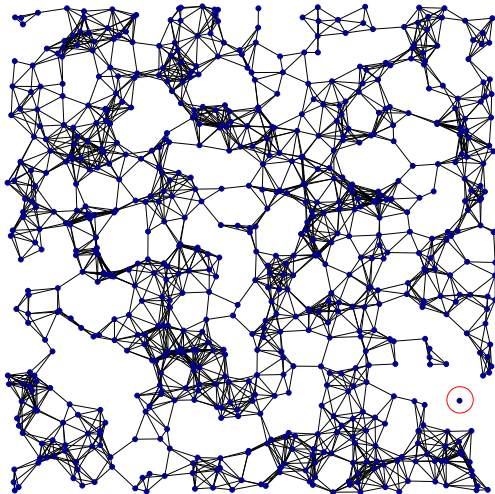
Random process:  $0 \leq r \leq \sqrt{2}$



$$r \sim \sqrt{C/n}$$

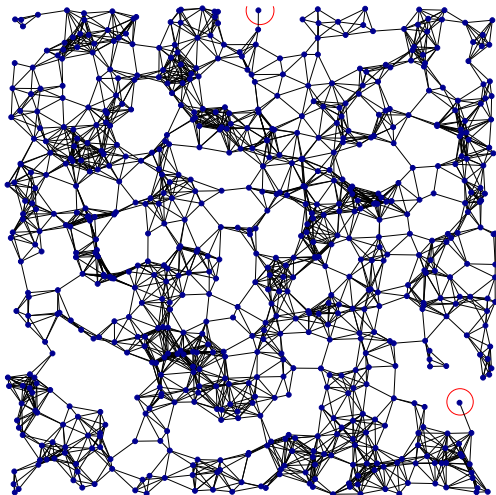
giant component!

Random process:  $0 \leq r \leq \sqrt{2}$



still disconnected!

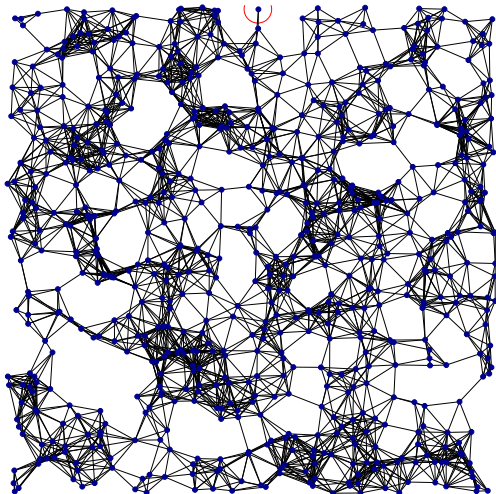
Random process:  $0 \leq r \leq \sqrt{2}$



connected  
=  
no isolated vertices  
(a.a.s.)

$$r = \sqrt{\frac{\log n + O(1)}{\pi n}}$$

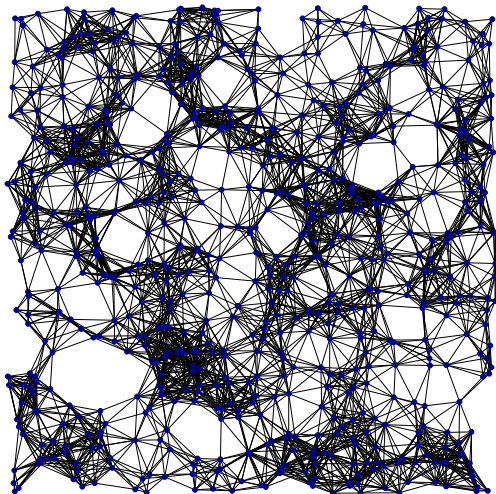
Random process:  $0 \leq r \leq \sqrt{2}$



2-connected  
=  
no deg. 1 vertices  
(a.a.s.)

$$r = \sqrt{\frac{\log n + \log \log n + O(1)}{\pi n}}$$

Random process:  $0 \leq r \leq \sqrt{2}$

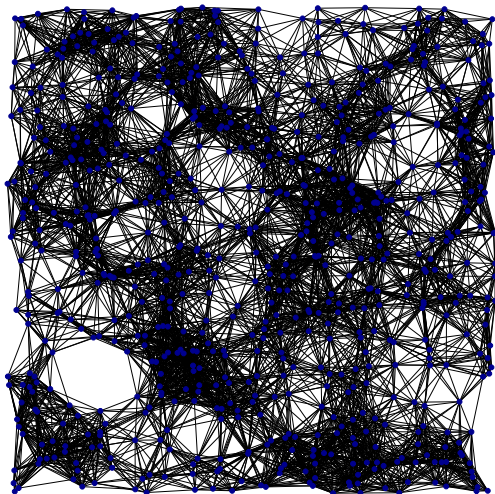


higher connectivity





Random process:  $0 \leq r \leq \sqrt{2}$



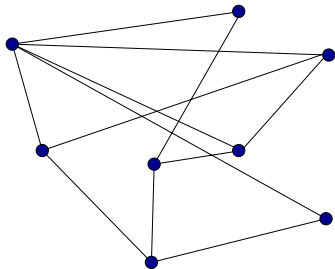
still large diameter:

$$\Theta(1/r)$$

bad expansion



## What about hamilton cycles?

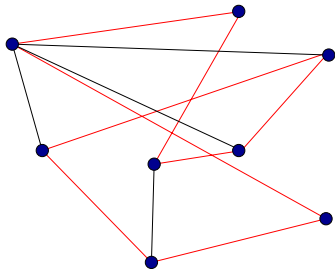


Necessary conditions:  $\min. \text{deg.} \geq 2$ , 2-connectivity

Are they sufficient for the RGG?



## What about hamilton cycles?



Necessary conditions:  $\text{min. deg.} \geq 2$ , 2-connectivity

Are they sufficient for the RGG?



## Earlier results

### Thm (Petit 2001)

The RGG with  $r = \sqrt{\omega(\log n)/n}$  has a.a.s. a Hamilton cycle.

### Thm (Díaz, Mitsche & P.G. 2007)

For any  $\epsilon > 0$ , the RGG with  $r \geq (1 + \epsilon)\sqrt{\frac{\log n}{\pi n}}$  has a.a.s. a Hamilton cycle.

(extension to general  $\ell_p$  norm)



## Recent results

Thm (Balogh, Bollobás, Krivelevich, Müller, P.G., Walters & Wormald 2010)

In the RGG process:

Hamiltonian  $\iff$  min. deg.  $\geq 2$  (a.a.s.)  
(extension to general dimension and  $\ell_p$  norm)

Thm (Balogh, Bollobás & Walters 2010)

Weaker analogue for the  $k$ -Nearest Neighbour Graph.

Thm (Krivelevich & Müller 2010)

Pancyclic  $\iff$  min. deg.  $\geq 2$  (a.a.s.)



## Recent results

Thm (Balogh, Bollobás, Krivelevich, Müller, P.G., Walters & Wormald 2010)

In the RGG process:

Hamiltonian  $\iff$  min. deg.  $\geq 2$  (a.a.s.)

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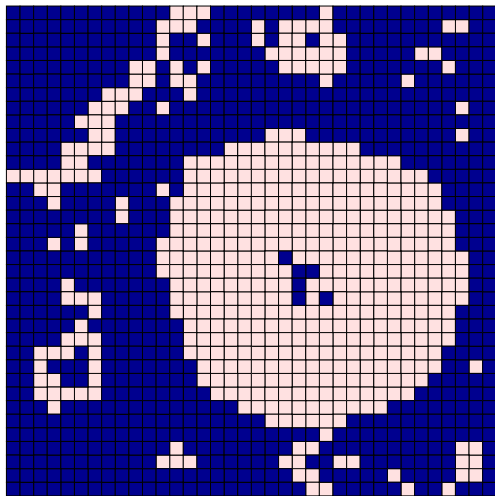
Thm (Müller, P.G. & Wormald 2010)

$k/2$  disjoint Hamilton cycles  $\iff$  min. deg.  $\geq k$  (a.a.s.)

(extension to general dimension and  $\ell_p$  norm)



# Main ideas: tessellation



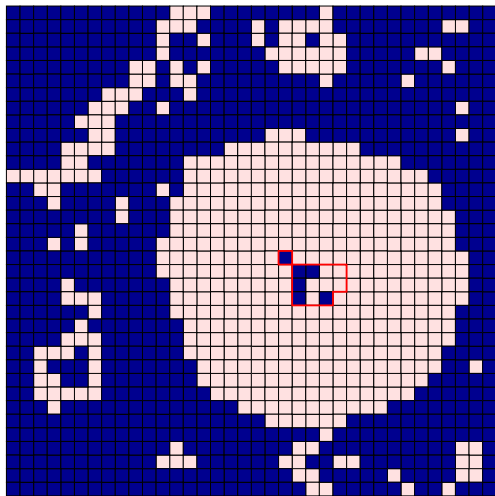
Set  $r = \sqrt{\frac{\log n}{\pi n}}$   
(not 2-connected)

□  $\downarrow \delta r$

■ dense ( $\geq M$  points)

□ sparse ( $< M$  points)

# Main ideas: tessellation

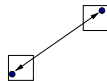


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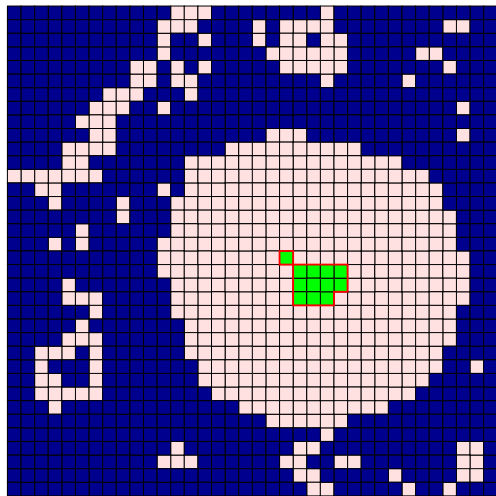
□ sparse ( $< M$  points)



dist.  $\leq r$



# Main ideas: tessellation

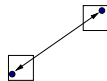


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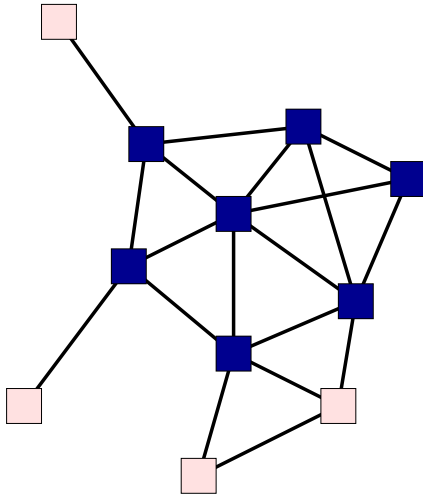


dist.  $\leq r$

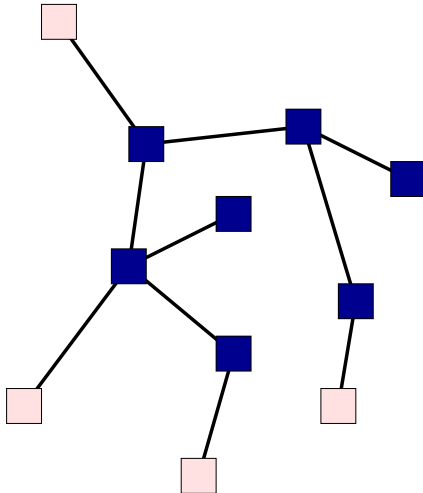


bad cells

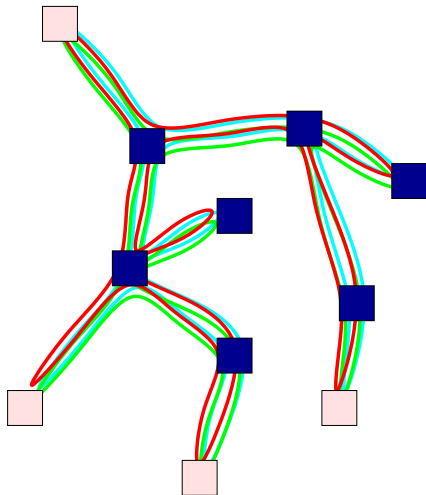
# Main ideas: large scale template



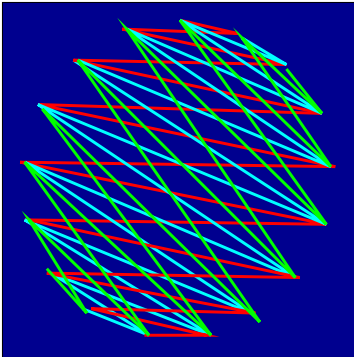
# Main ideas: large scale template



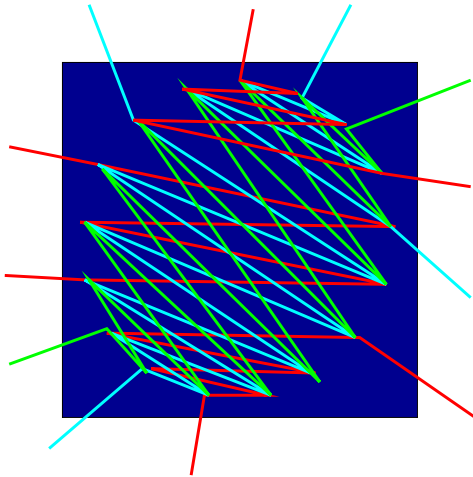
# Main ideas: large scale template



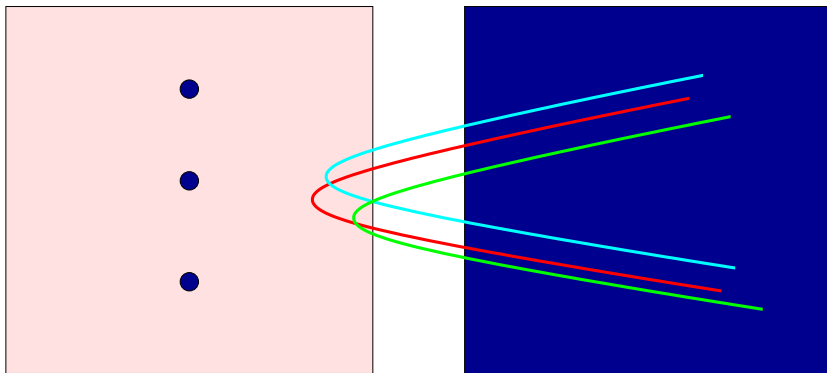
# Main ideas: rerouting I



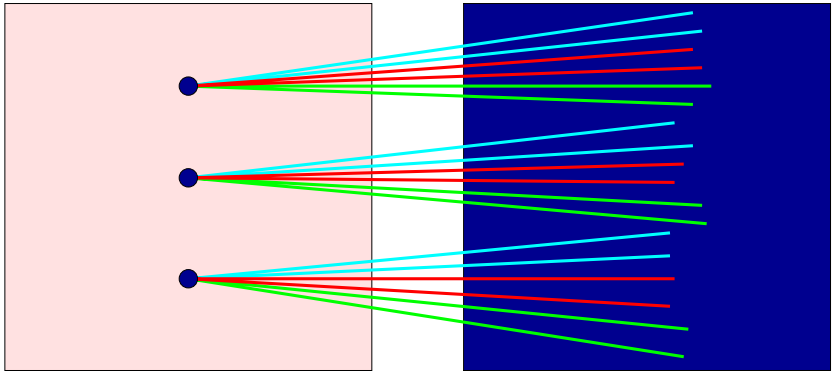
# Main ideas: rerouting I



# Main ideas: rerouting II

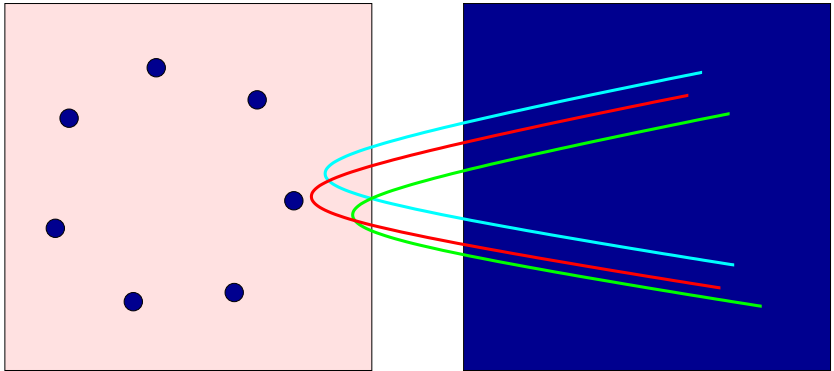


# Main ideas: rerouting II

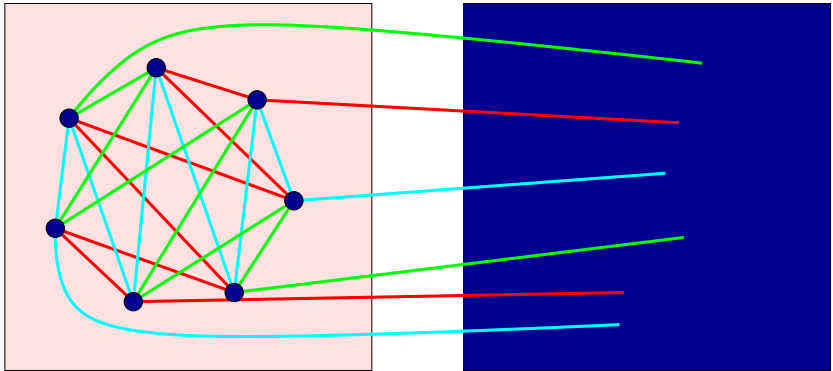




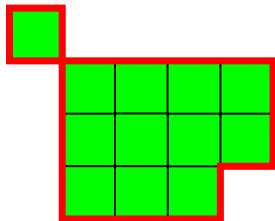
# Main ideas: rerouting II



# Main ideas: rerouting II



# Main ideas: rerouting III



a bit harder!



# Open question

Are there always  $\lfloor \frac{\delta(RGG)}{2} \rfloor$  edge disjoint Hamilton cycles?



thank you!



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